#### MATENVMED

#### Δράση 3.2 ΥΛΟΠΟΙΗΣΗ ΣΕ FPGAs ΚΑΙ ΠΟΛΥΠΥΡΗΝΑ ΣΥΣΤΗΜΑΤΑ (MULTICORES / MANYCORES)

#### Νίκος Μπέλλας Τμήμα Ηλεκτρολόγων Μηχανικών και Μηχανικών Η/Υ Πανεπιστήμιο Θεσσαλίας



#### Deliverable 3.2



- Select and study computationally demanding kernels of WP2 that can potentially be accelerated when mapped to a massively parallel platform (GPUs, FPGAs).
- Design and implement appropriate tools or use third-party tools to help with mapping these kernels to such platforms (compilers, CAD tools)
- Map the kernels to FPGAs and GPUs
- Performance analysis and comparison wrt. to CPU implementation

# What is an FPGA?



- *Field Programmable Gate Array (FPGA)* is the best known example of Reconfigurable Logic
- Hardware can be modified post chip fabrication
- Tailor the Hardware to the application
  - Fixed logic processors (CPUs/GPUs) only modify their software (via programming)
- FPGAs can offer superior performance, performance/power, or performance/cost compared to CPUs and GPUs.





#### FPGA architecture



- A generic island-style FPGA fabric
- Configurable Logic Blocks (CLB) and Programmable Switch Matrices (PSM)
- **Bitstream** configures functionality of each CLB and interconnection between logic blocks



# FPGA discussion



- Advantages
  - Hardware tailored to application: potential for (near) optimal performance for a given application
  - Various forms of parallelisms can be exploited
- Disadvantages
  - Programmable mainly at the hardware level using Hardware Description Languages (BUT, this can change)
  - Lower clock frequency (< 300 MHz) compared to CPUs (~ 3GHz) and GPUs (~1.5 GHz)

# OpenCL for Heterogeneous Systems

- OpenCL (Open Computing Language) : A unified programming model aims at letting a programmer write a portable program once and deploy it on any heterogeneous system with CPUs and GPUs.
- Became an important industry standard after release due to substantial industry support.



# **OpenCL Platform Model**



One host and one or more Compute Devices (CD) Each CD consists of one or more Compute Units (CU) Each CU is further divided into one or more Processing Elements (PE) MATENVMED Final Review

# **OpenCL Kernel Execution Geometry**

- OpenCL defines a geometric partitioning of grid of computations
- Grid consists of N dimensional space of work-groups
- Each work-group consists of N dimensional space of work-items.



# **OpenCL Simple Example**

- OpenCL kernel describes the computation of a workitem
  - Finest parallelism granularity
- e.g. add two integer vectors (N=1)  $\frac{R_{L}}{L}$

Run-time call

Used to differentiate execution for each work-item

```
kernel void vadd(
void add(int* a,
                                                                  _global int* a,
     int* b,
                                                                  _global int* b,
     int^* c) {
                                                                __global int* c) {
for (int idx=0; idx<sizeof(a); idx++)
                                                             int idx= get_global_id(0);
 c[idx] = a[idx] + b[idx];
                                                             c[idx] = a[idx] + b[idx];
                C code
                                                                  OpenCL kernel code
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                                                                                                    9
```

# Silicon OpenCL

- Silicon-OpenCL "SOpenCL".
- A tool flow to convert an unmodified OpenCL application into a SoC design with HW/SW components.
- A template-based hardware accelerator generation.
- Decouple data movement and computations.



Input

data

# **OpenCL Implementation**



- Our plan was to use the same code base (e.g. OpenCL) to explore different architectures
  - OpenCL used for multicore CPU, GPU, FPGA (SOpenCL)
- Fast exploration based on area, performance and power requirements

### Monte Carlo random walks

All random walks from all points are independent.

Double precision arithmetic with non-constant number of iterations

Compute-bound algorithm



$$do\{rad = rand\_uniform(seed) * d; \\ while(\frac{4 * rad}{(d * d)} * log(\frac{d}{rad}) < \frac{4 * rand\_uniform()}{e * d}); \\ 14/11/2015 \qquad MATENVMED Final Review$$

# **CPU - GPU implementation**

- Intel Core i7-4820K CPU clocked at 3.70GHz
  - 8 threads
  - Each thread performs a random walk
- GeForce GTX 680 GPU (Kepler architecture)
  - 3 Teraflops peak performance, 288 Watts power dissipation (TDP). 1536 cores. 3.5 billion transistors.
  - N\*768 threads (N is the number of points in the grid)
  - Each thread performs a random walk



# FPGA implementation (I)



- Good match for FPGA technology:
  - MC algorithm massively parallel
  - Independent multiple path traversal from multiple points
- Poor match for FPGA technology:
  - Double precision (DP) Trigonometric, Log, Additions, Multiplications functions for each walk
  - DP arithmetic takes up a lot of area and is slow

$$\begin{aligned} &do \{ \\ &rad = rand\_uniform(seed) * d; \\ &\}while(\frac{4 * rad}{(d * d)} * log(\frac{d}{rad}) < \frac{4 * rand\_uniform()}{e * d}); \end{aligned}$$

# FPGA implementation (II)



- Use SOpenCL and Xilinx Vivado HLS (High Level Synthesis) tools to automatically generate MC hardware accelerators
- Multiple accelerators with different performance vs. area characteristics can be generated very fast
- Best approach is when one accelerator traverses all walks from as many points as possible
- Target clock 250 MHz

# **FPGA** architecture





Parameterizable Architecture in Number of Accelerators Automatically generated after tools in place (SOpenCL and Vivado HLS)

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# Virtex-7 FPGA implementation



Desktop-based High Performance Accelerator 5 MC accelerators shown





#### Performance





## Power and Power Efficiency





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Multicore CPU

1.0

5,0

10,0

15,0

0,0

25,0

20,0

#### MATENVMED-sponsored Publications (Δράση 3.2)



- Muhsen Owaida, Christos D. Antonopoulos, Nikolaos Bellas. A Grammar Induction Method for Clustering of Operations in Complex FPGA Designs. *IEEE Symposium on Field-Programmable Custom Computing Machines (FCCM)*, May 11-13, 2014, Boston, MA.
- K. Krommydas, W.C. Feng, M. Owaida, C.D. Antonopoulos, and N. Bellas. On the Portability of the OpenCL Dwarfs on Fixed and Reconfigurable Parallel Platforms. *IEEE International Conference on Application-Specific Systems, Architectures and Processors (ASAP).* June 18-20, 2014, Zurich, Switzerland. (Nominated for Best Paper Award)
- Christos Antonopoulos. SOpenCL: An Infrastructure for Transparently Integrating FPGAs in Heterogeneous, Accelerator-Based Systems. 6<sup>th</sup> Conference on Numerical Analysis (NumAn). September 2-5, 2014. Chania.
- Panagiotis Skribonis, George Zindros, Ioannis Parnassos, Muhsen Owaida, Nikolaos Bellas, Paolo Ienne. Exploring Automatically Generated Platforms in High Performance FPGAs. *Parallel FPGA* (*ParaFPGA*). September 1-4, 2015. Edinburgh, UK
- K. Krommydas, W.C. Feng, C.D. Antonopoulos, and N. Bellas. On the Characterization of OpenCL Dwarfs on Fixed and Reconfigurable Platforms. *Journal of Signal Processing Systems (JSSP)*. US Springer. October 2015.

### Conclusion



- Graphics Processor Units (GPUs) are ideally positioned to run massively parallel applications
- FPGAs the same, but Double Precision Arithmetic is not their strong point
- Better Hardware Libraries for DP arithmetic are needed.
- Manycore parallel architectures mature to be used in framework-based computing